Before spring break notes

Prof Bill, Feb 2020

Before we head to the beach, I want to introduce balanced trees.

One page per tree. We'll go *deep* after Spring break.

The goal of balanced trees is to eliminate the destructive case where a BST turns into a linked list. In other words, we want the worst case performance of our balanced trees to be O(log n).

Each page starts with a Wikipedia link as an introduction. Then, I add some Bill-blather. Pages end with an animation link, so we can play. (yeah!)

Pages:

- 1. B-trees, 2-3-4 trees
- 2. Red-black trees
- 3. AVL trees
- 4. Scapegoat trees

thanks...yow, bill

1. B-trees, 2-3-4 trees

I like Wikipedia here, actually:

- → B-tree, en.wikipedia.org/wiki/B-tree
- → 2-3-4 tree, en.wikipedia.org/wiki/2%E2%80%933%E2%80%934_tree

B-trees

Important: "Not to be confused with Binary tree"

Definition:

In computer science, a B-tree is **a self-balancing tree** data structure that maintains sorted data and allows searches, sequential access, insertions, and deletions in **logarithmic time**. The B-tree generalizes the binary search tree, allowing for **nodes with more than two children**. Unlike other self-balancing binary search trees, the B-tree is well suited for storage systems that read and write relatively large blocks of data, such as discs. It is commonly used in databases and file systems.

2-3-4 trees

We'll focus on 2-3-4 trees, which are one flavor of B-tree. Definition:

In computer science, a 2–3–4 tree (also called a 2–4 tree) is a self-balancing data structure...where every node with children (internal node) has either two, three, or four child nodes:

- a 2-node has one data element, and if internal has two child nodes
- a 3-node has two data elements, and if internal has three child nodes
- a 4-node has three data elements, and if internal has four child nodes



Properties:

- Every node (leaf or internal) is a 2-node, 3-node or a 4-node, and holds one, two, or three data elements, respectively.
- > All leaves are at the same depth (the bottom level).
- ➤ All data is kept in sorted order.

Since our tree is always balanced, then search, insert, and delete are all O(log n)!

2–3–4 trees are an isometry of red–black trees, meaning that they are equivalent data structures. Keep reading...

Maybe the best animation of all (select "B Tree"): www.cs.usfca.edu/~galles/visualization/Algorithms.html

2. Red-black trees

Wikipedia goodness: en.wikipedia.org/wiki/Red%E2%80%93black_tree

The red-black rules are:

- **□** Each node is either red or black.
- □ The root is black.
- □ All leaves (NIL) are black.
- □ If a node is red, then both its children are black.
- □ Every path from a given node to any of its descendant NIL nodes contains the same number of black nodes.



Great animation (select "Red-Black Trees"): www.cs.usfca.edu/~galles/visualization/Algorithms.html

3. AVL trees

Wikipedia: https://en.wikipedia.org/wiki/AVL_tree

Definition.

In computer science, an AVL tree (named after inventors Adelson-Velsky and Landis) is a self-balancing binary search tree. It was the first such data structure to be invented.[2] In an AVL tree, the heights of the two child subtrees of any node differ by at most one; if at any time they differ by more than one, rebalancing is done to restore this property.

AVL is notoriously complex/difficult to code. You can see...Wikipedia hides the huge code chunks.

We'll learn the various rotations to make AVL work.

This is an OK start. (hint: we will do better) www.tutorialspoint.com/data_structures_algorithms/avl_tree_algorithm.htm

4. Scapegoat trees

Wikipedia: en.wikipedia.org/wiki/Scapegoat_tree

Read Morin 8 Scapegoat trees: <u>opendatastructures.org/ods-java/8_Scapegoat_Trees.html</u>

Quick facts: 1) The Wikipedia page for scapegoat trees is weak(ipedia), 2) There is no animation for scapegoat trees, and 3) I've never used them.

/* What do these facts tell you? */