

# Red Black Tree insertion

Some pseudocode:

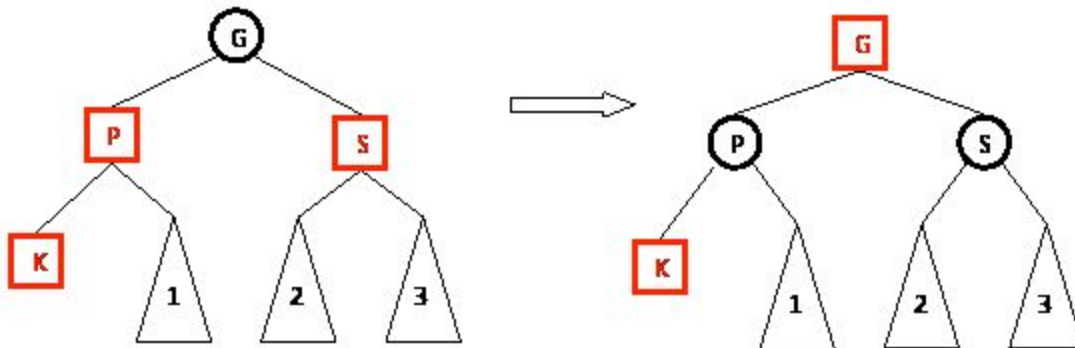
```
insert( K key)
  n = create red node( key)
  if empty tree
    root = n
    change n color to black
  else
    do BST insert of n as leaf
    if parent of n is red
      // new node and its parent are both red = must fix this!
      if uncle of n is red, then recolor
      else rotate
```

## Recoloring

In the figure below... a new node (K) is added. It's parent (P) is red causing a red-red violation. If the uncle (S) is red, then recolor in two steps:

1. Make the grandparent (G) red, and
2. Color its children (P and S) black.

This resolves the red-red conflict AND maintains equal black-height.



Important - making the red grandparent (G) may cause a conflict above us. Apply the same recursively to grandparent (G).

## Rotation

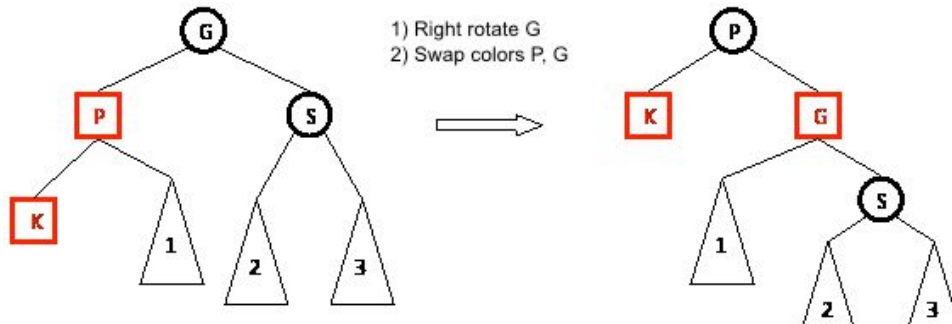
If the new node's uncle is a black node, then rotation is required.

There are 4 rotations cases (similar to AVL). They're on the next back (the back).

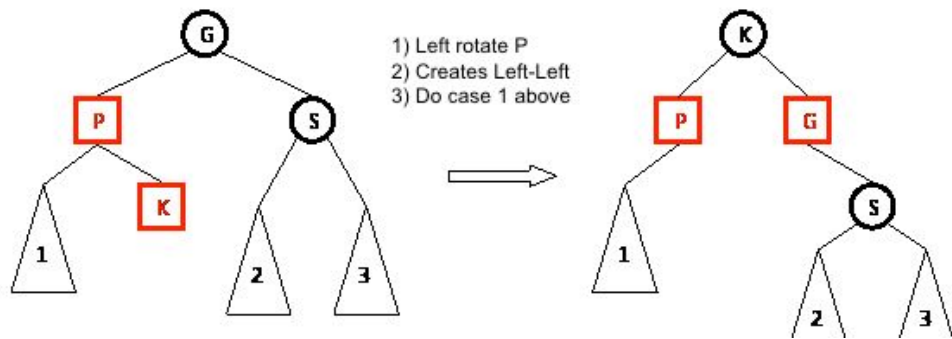
Source: [pages.cs.wisc.edu/~paton/readings/Red-Black-Trees/](http://pages.cs.wisc.edu/~paton/readings/Red-Black-Trees/)

**4 Rotation cases** - when new node (K) has a red parent (P) and a black uncle (S).

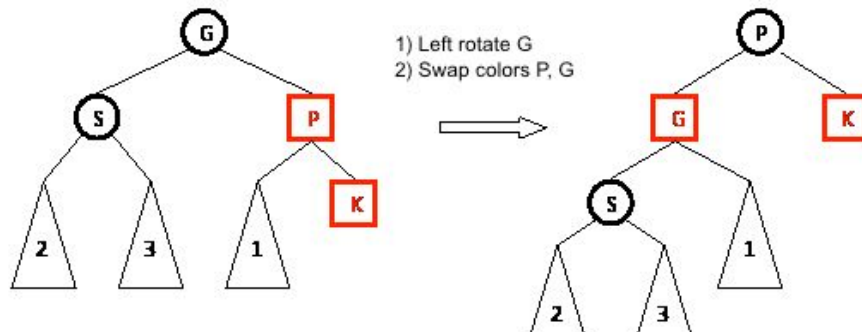
Case 1: Left-Left (P is parent, K is new (Key) node)



Case 2: Left-Right



Case 3: Right-Right



Case 4: Right-Left

